Typal Heterogeneous Equality Types

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The usual homogeneous form of equality type in Martin-Löf Type Theory contains identifications between elements of the same type. By contrast, the heterogeneous form of equality contains identifications between elements of possibly different types. This short note introduces a simple set of axioms for such types. The axioms are shown to be equivalent to the combination of systematic elimination rules for both forms of equality, albeit with *typal* (also known as “propositional”) computation properties, together with Streicher’s Axiom K, or equivalently, the principle of uniqueness of identity proofs.

CCS Concepts: • **Theory of computation** → **Type theory**;

Additional Key Words and Phrases: Type theory, equality types, heterogeneous equality

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1 INTRODUCTION

Equality types in the intensional version of Martin-Löf Type Theory (see, for example, Reference [12, Section 8.1]) are traditionally formulated in terms of an introduction rule (reflexivity) together with a rule for eliminating proofs of equality and a rule describing how elimination computes when it meets a reflexivity proof. Some recent work [4, 7] on models of Homotopy Type Theory [17] uses a formulation of equality types that differs from this in two respects. First, the elimination operation is replaced by the combination of a simple operation for transporting elements along proofs of equality, together with an axiom asserting contractibility of singleton types. Second, the analogue of the computation rule for the eliminator, namely, that transporting along a reflexivity proof does nothing, is weakened from a judgemental equality to the existence of an element of the corresponding equality type; see Reference [9] and Figure 2 in Reference [4]. This formulation is sometimes called a “propositional” equality type [18], but here I will follow Shulman [14, Section 1.6] for the reasons given there and refer to *typal* equality types. Although these changes to the formulation of equality types affect computation, it seems that they do not change what is provable (see Reference [5], for example) and they make it easier to construct models. Furthermore, they can lead to surprising simplifications. For example, Lumsdaine [private communication] has observed that the computation rule is superfluous (for elimination, but the observation also holds for transport): If a proto-identity type has a transport operation lacking its

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Axioms for typal heterogeneous equality satisfying Axiom K

```

postulate
  _≡≡_ : ∀{l}{A B : Set l} → A → B → Set l

- the derived homogeneous equality
_≡_   : ∀{l}{A : Set l} → A → A → Set l
x ≡ y = x ≡≡ y

postulate
  rfl : ∀{l}{A : Set l} (x : A) → x ≡ x
  ctr : ∀{l}{A B : Set l}{x : A}{y : B}(e : x ≡≡ y) → rfl x ≡≡ e
  eqt : ∀{l}{A B : Set l}{x : A}{y : B} → x ≡≡ y → A ≡ B
  tpt : ∀{l m n}{A : Set l}{B : A → Set m}(C : (x : A) → B x → Set n)
        {x x' : A}{y : B x}{y' : B x'} → x ≡ x' → y ≡≡ y' → C x y → C x' y'

Axioms for Σ-types with surjective pairing

postulate
  Σ : ∀{l m}(A : Set l)(B : A → Set m) → Set (l ⊔ m)

module _ {l m}{A : Set l}{B : A → Set m} where
  postulate
    _,_ : (x : A) → B x → Σ A B
    fst : Σ A B → A
    snd : (z : Σ A B) → B (fst z)
    fpr : (x : A)(y : B x) → fst(x , y) ≡ x
    spr : (x : A)(y : B x) → snd(x , y) ≡≡ y
    eta : (z : Σ A B) → (fst z , snd z) ≡ z

- concrete syntax for Σ-types
syntax Σ A (λ x → B) = Σ x : A , B

```

Fig. 1. The Axioms.

typal computation property, then the operation can be corrected to a new one that does have the computation property (see Lemma 2.1 and the Appendix).

The above remarks apply to the usual, homogeneous notion of equality in which elements of the same type are compared. The purpose of this article is to give an analogous treatment of *heterogeneous* equality [2, 11] in the presence of Σ -types and the Axiom K of Streicher [15, Section 1.2]. Since Axiom K is not compatible with the Univalence Principle of Homotopy Type Theory [17, Example 3.1.9], the focus here is on the simpler (but still useful!) world of zero-dimensional type theory. We will see that the axioms in Figure 1 capture homogeneous and heterogeneous equality satisfying their usual dependent elimination and (typal) computation properties and Axiom K, and Σ -types with their usual dependent elimination and (typal) computation properties. It seems necessary to include Σ -types to get Lumsdaine's result mentioned above (see Remark 2.5); the axioms we give for such types are standard, except that the equality property of dependent second projection (*spr*) is simplified by the use of heterogeneous rather than homogeneous equality. The axioms in the figure are pleasingly simple compared to the usual formulation in terms of elimination and computation properties, and may aid finding new models of heterogeneous equality types.

The implementation of intensional Martin-Löf Type Theory provided by Agda 2.6 [1] is used to state the axioms and develop their properties. More precisely, we just make use of Agda's implementation of a countably infinite, non-cumulative hierarchy of universes `Set l`, where `l` ranges over a type `Level` of universe levels whose closed normal forms are in bijection with the natural numbers. The universes are closed under dependent function types (written in Agda as $(x : A) \rightarrow B$) and inductive types. The use of a whole hierarchy of universes is necessary; for example, the function `eqt` in Figure 1 takes a heterogeneous equality type $x \equiv y$ in universe `Set l` and produces a homogeneous one $A \equiv B$ in the universe one level up, which is denoted `Set(lsuc l)` in Agda. We also use Agda's notation for infix and for implicit arguments. For example, the function `_≡_` in Figure 1 takes five arguments, the first three of which are implicit and the last two of which are infix. In particular, Agda's ability to infer the values of implicit arguments (or of unspecified explicit arguments, which are denoted by an underscore, `_`) is used quite aggressively in what follows, to be able to see the wood from the trees.

Although the code in this article has been checked by Agda, some parts of it that are not essential for understanding the development have been hidden; the complete (non-literate) Agda code can be found at <https://doi.org/10.17863/CAM.47902>.

2 THE AXIOMS AND THEIR PROPERTIES

Figure 1 postulates a family of types `_≡_` in all universes, together with some operations on them that together capture a typal version of *heterogeneous* equality. Heterogeneous equality types were introduced by McBride [11, Section 5.1.3] under the name of "John Major equality." Unlike ordinary, homogeneous equality types, such a type $x \equiv y$ relates elements x and y of possibly different types, A and B say. The intention is that elements of type $x \equiv y$ denote proofs that not only are x and y equal, but so also are their types A and B . The figure defines homogeneous equality `_≡_` as the special case of `_≡_` when the types of the two arguments are already known to be the same. Axiom `rfl` says that `≡` is reflexive. Axiom `ctr` is a heterogeneous version of the contractibility property of singleton types (cf. center in Figure 2 of [4]). Axiom `eqt` says that heterogeneously equal things have (homogeneously) equal types. Axiom `tpt` is a form of the transport property of equality (cf. T in Figure 2 of Reference [4]) involving both homogeneous and heterogeneous equalities. Finally, `Σ`, `_`, `fst`, `snd`, `fpr`, `spr` and `eta` axiomatize dependent product types satisfying surjective pairing.

We begin with some simple lemmas establishing the basics of equational logic for `≡`, namely, chain-reasoning using reflexivity (already an axiom), symmetry, transitivity, and congruence properties. These are given in Figure 2.

The axioms in Figure 1 are notably lacking a "regularity" property for `tpt`, that is, a proof of type `tpt (rfl x) (rfl y) z ≡ z`. But such a thing is needed if we are to derive the expected elimination and (typal) computation rules for `_≡_` and `_≡_`. To get those, one can define a "corrected" form of transport that has this regularity property, using a simplified version of a trick due to Peter Lumsdaine [unpublished]. In fact, it is enough to produce a function coercing proofs of equality of types $e : A \equiv B$ into functions `coe e : A → B` and which satisfies the heterogeneous regularity property that `coe e x ≡ x` (so that, given how we define `≡` in terms of `≡`, the usual form of regularity, `coe (rfl A) x ≡ x`, is just the special case of this when e is `rfl A`).

LEMMA 2.1. *The axioms in Figure 1 imply the existence of a coercion function*

`coe : ∀{l}{A B : Set l} → A ≡ B → A → B`

satisfying a heterogeneous regularity property:

`coeIsRegular : ∀{l}{A B : Set l}(e : A ≡ B)(x : A) → coe e x ≡ x`

```

symm :  $\forall\{l\}\{A B : \text{Set } l\}\{x : A\}\{y : B\} \rightarrow x \equiv y \rightarrow y \equiv x$ 
symm e = tpt ( $\lambda \_ y \rightarrow y \equiv \_$ ) (eqt e) e (rfl  $\_$ )

proof_ :  $\forall\{l\}\{A B : \text{Set } l\}\{x : A\}\{y : B\} \rightarrow x \equiv y \rightarrow x \equiv y$ 
proof p = p

_ $\equiv$ [_]_ :  $\forall\{l\}\{A B C : \text{Set } l\}\{x : A\}\{y : B\}\{z : C\} \rightarrow$ 
     $x \equiv y \rightarrow y \equiv z \rightarrow x \equiv z$ 
x  $\equiv$  [ e ] f = tpt ( $\lambda \_ z \rightarrow x \equiv z$ ) (eqt f) f e

_qed :  $\forall\{l\}\{A : \text{Set } l\}\{x : A\} \rightarrow x \equiv x$ 
x qed = rfl x

cong :  $\forall\{l m\}\{A : \text{Set } l\}\{B : A \rightarrow \text{Set } m\}\{f : (x : A) \rightarrow B x\}\{x y : A\} \rightarrow$ 
     $x \equiv y \rightarrow f x \equiv f y$ 
cong f {x} e = tpt ( $\lambda \_ z \rightarrow f x \equiv f z$ ) e e (rfl (f x))

cong2 :  $\forall\{l m n\}\{A : \text{Set } l\}\{B : A \rightarrow \text{Set } m\}\{C : (x : A) \rightarrow B x \rightarrow \text{Set } n\}$ 
     $(f : (x : A) (y : B x) \rightarrow C x y)\{x x' : A\}\{y : B x\}\{y' : B x'\} \rightarrow$ 
     $x \equiv x' \rightarrow y \equiv y' \rightarrow f x y \equiv f x' y'$ 
cong2 f {x} {_-} {y} e e' = tpt ( $\lambda x' y' \rightarrow f x y \equiv f x' y'$ ) e e' (rfl (f x y))

cong3 :  $\forall\{k l m n\}\{A : \text{Set } k\}\{B : A \rightarrow \text{Set } l\}\{C : (x : A) \rightarrow B x \rightarrow \text{Set } m\}$ 
     $\{D : (x : A) (y : B x) \rightarrow C x y \rightarrow \text{Set } n\}$ 
     $(f : (x : A) (y : B x) (z : C x y) \rightarrow D x y z)$ 
     $\{x x' : A\}\{y : B x\}\{y' : B x'\}\{z : C x y\}\{z' : C x' y'\} \rightarrow$ 
     $x \equiv x' \rightarrow y \equiv y' \rightarrow z \equiv z' \rightarrow f x y z \equiv f x' y' z'$ 
cong3 f {x} {_-} {y} {_-} {z} e e' =
    tpt ( $\lambda x' y' \rightarrow \forall\{z'\} \rightarrow z \equiv z' \rightarrow f x y z \equiv f x' y' z'$ ) e e' (cong2 (f x) (rfl y))

```

Fig. 2. Equational reasoning for heterogeneous equality.

PROOF. First, we define the type of functions that are injective with respect to $\equiv$ and note that the identity function is one such:

```

Inj :  $\forall\{l\}\{A B : \text{Set } l\} \rightarrow \text{Set } l$ 
Inj A B =  $\sum f : (A \rightarrow B), \forall\{x y\} \rightarrow f x \equiv f y \rightarrow x \equiv y$ 
id :  $\forall\{l\}\{A : \text{Set } l\} \rightarrow A \rightarrow A$ 
id x = x

```

```

idInj :  $\forall\{l\}\{A : \text{Set } l\} \rightarrow \text{Inj } A A$ 
idInj _ = (id, id)

```

Next, we use **tpt** to define a function coercing equalities into injective functions:

```

icoe :  $\forall\{l\}\{A B : \text{Set } l\} \rightarrow A \equiv B \rightarrow \text{Inj } A B$ 
icoe {l} {A} e = tpt ( $\lambda \_ C \rightarrow \text{Inj } A C$ ) (rfl (Set l)) e (idInj A)

```

The injectiveness of **icoe** e is used as follows. Applying the operation **tpt** to the type family

```

fsticoe :  $\forall\{l\}\{A : \text{Set } l\}\{x : A\}\{B : \text{Set } l\}\{e : A \equiv B\} \rightarrow \text{Set } l$ 
fsticoe x B e =  $\sum y : B, (\text{fst } (\text{icoe } (\text{rfl } B)) y \equiv \text{fst } (\text{icoe } e) x)$ 

```

we can transport the element $(e, \text{rfl}(\text{fst}(\text{icoe}(\text{rfl } A)) x))$ of type $\text{fsticoe } x A (\text{rfl } A)$ along $e : A \equiv B$ and $\text{ctr } e : \text{rfl } A \equiv e$ to give an element of type $\text{fsticoe } x B e$. The first projection of this element gives the value of the desired coercion along e at x :

$\text{coe } e x = \text{fst}(\text{tpt}(\text{fsticoe } x) e (\text{ctr } e) (x, \text{rfl } _))$

and its second projection can be used along with the injectiveness property of icoe to get the regularity property of this coercion:

```
coeIsRegular {_} {A} e x = tpt (λ _ e' → coe e' x ≡ x) e (ctr e) coerfl
  where
    coerfl : coe (rfl A) x ≡ x
    coerfl = snd (icoe (rfl A)) (snd (
      tpt (fsticoe x) (rfl A) (ctr (rfl A)) (x, rfl _)))
```

□

An immediate corollary is that the axioms imply the *uniqueness of identity proofs* (UIP) and hence Streicher's Axiom K [15]. (We will see in Section 3 that in fact it is only the tpt function that contains an implicit use of Axiom K.)

THEOREM 2.2 (UIP AND AXIOM K). *The axioms in Figure 1 imply that $\equiv$ satisfies*

uip : $\forall \{l\} \{A : \text{Set } l\} \{x y : A\} (e e' : x \equiv y) \rightarrow e \equiv e'$

axiomK : $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : x \equiv x \rightarrow \text{Set } m) (p : P (\text{rfl } x)) \rightarrow \forall e \rightarrow P e$

axiomKComp : $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : x \equiv x \rightarrow \text{Set } m) (p : P (\text{rfl } x)) \rightarrow \text{axiomK } P p (\text{rfl } x) \equiv p$

PROOF. Using the functions from Figure 2 and Lemma 2.1, we have

uip $e e' = \text{tpt}(\lambda _ e'' \rightarrow e'' \equiv e') e (\text{ctr } e) (\text{ctr } e')$

axiomK $P p e = \text{coe}(\text{cong}_2(\lambda _ \rightarrow P) (\text{rfl } p) (\text{ctr } e)) p$

axiomKComp $P p = \text{coeIsRegular}(\text{cong}_2(\lambda _ \rightarrow P) (\text{rfl } p) (\text{ctr}(\text{rfl } _))) p$

□

The elimination and computation properties of $\equiv$ and $\equiv\equiv$ then follow:

THEOREM 2.3 (ELIMINATION AND TYPAL COMPUTATION PROPERTIES). *The axioms in Figure 1 imply that $\equiv$ has the usual elimination and (typal) computation properties of homogeneous equality (in the form suggested by Reference [13])*

$\equiv\text{Elim}$: $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : (y : A) \rightarrow x \equiv y \rightarrow \text{Set } m) (p : P x (\text{rfl } x)) (y : A) (e : x \equiv y) \rightarrow P y e$

$\equiv\text{Comp}$: $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : (y : A) \rightarrow x \equiv y \rightarrow \text{Set } m) (p : P x (\text{rfl } x)) \rightarrow \equiv\text{Elim } P p x (\text{rfl } x) \equiv p$

The axioms also imply that $\equiv\equiv$ has the elimination and (typal) computation properties of heterogeneous equality described by McBride [11, Section 5.1.3]

$\equiv\equiv\text{Elim}$: $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : (B : \text{Set } l) (y : B) \rightarrow x \equiv\equiv y \rightarrow \text{Set } m) (p : P A x (\text{rfl } x)) (B : \text{Set } l) (y : B) (e : x \equiv\equiv y) \rightarrow P B y e$

$\equiv\equiv\text{Comp}$: $\forall \{l m\} \{A : \text{Set } l\} \{x : A\} (P : (B : \text{Set } l) (y : B) \rightarrow x \equiv\equiv y \rightarrow \text{Set } m) (p : P A x (\text{rfl } x)) \rightarrow \equiv\equiv\text{Elim } P p A x (\text{rfl } x) \equiv p$

PROOF. Using the functions from Figure 2 and Lemma 2.1, we have

```

≡Elim P p _ e   = coe (cong2 P e (ctr e)) p
≡Comp P         = coeIsRegular (cong2 P (rfl _) (ctr (rfl _)))
≡≡Elim P p _ _ e = coe (cong3 P (eqt e) e (ctr e)) p
≡≡Comp P        = coeIsRegular (cong3 P (eqt (rfl _)) (rfl _) (ctr (rfl _)))

```

□

Note that a corollary of the above two theorems is that $_ \equiv _$ is uniquely determined up to logical equivalence by the axioms in Figure 1. In other words, for any other such family of types $_ \equiv' _$, there are functions in either direction between $x \equiv y$ and $x \equiv' y$; and because of UIP, these are necessarily mutually inverse up to $\equiv$ (or $\equiv'$).

Remark 2.4. $\equiv\equiv\text{Elim}$ is the elimination form systematically derived [3] from $_ \equiv _$ and rfl , regarding them as the formation and introduction rules for an inductive type. As McBride [11, page 120] points out, $\equiv\equiv\text{Elim}$ is not very useful because of the way its motive P involves abstraction over an arbitrary type B . McBride goes on to give another, more useful form of elimination for $\equiv$, but in our setting where $\equiv$ is a special case of $\equiv\equiv$, that coincides with the eliminator $\equiv\text{Elim}$.

Remark 2.5 (The role of Σ -types). One of the strengths of machine-checked mathematics is that it aids the detection of logical dependency. Although we included the equations fpr , spr , and eta for Σ -types in Figure 1, they have not been used for the results so far, as may be verified by commenting them out from this literate Agda file and re-checking it up to this point.

So only the weak form of dependent product given by Σ , $_$, fst and snd in the figure is used to define the regular version of coercion in Lemma 2.1 and then prove Theorems 2.2 and 2.3. It would be nice if there was some way to define Σ , $_$, fst , and snd just using dependent function types and universes.

However, the extra equations fpr , spr , and eta for Σ are of course very natural. Let us record the fact that they enable one to define the usual elimination rule for dependent products, with a typical computation rule:

```

ΣElim      : ∀ {l m n} {A : Set l} {B : A → Set m} {C : Σ A B → Set n}
              (c : (x : A) (y : B x) → C (x , y)) (z : Σ A B) → C z
ΣElim C c z = coe (cong C (eta z)) (c (fst z) (snd z))

ΣComp      : ∀ {l m n} {A : Set l} {B : A → Set m} {C : Σ A B → Set n}
              (c : (x : A) (y : B x) → C (x , y)) (x : A) (y : B x) →
              ΣElim C c (x , y) ≡ c x y
ΣComp C c x y = let z = (x , y) in
proof
  coe (cong C (eta z)) (c (fst z) (snd z))
≡≡[ coeIsRegular _ _ ]
  c (fst z) (snd z)
≡≡[ cong2 c (fpr x y) (spr x y) ]
  c x y
qed

```

3 CONSISTENCY OF THE AXIOMS

We have seen that the axioms in Figure 1 suffice to define dependent products and both heterogeneous and homogeneous equality types with uniqueness of identity proofs, all satisfying the usual elimination properties, albeit with typal computation rules. Conversely, it is not hard to see that the elimination and computation rules in Theorem 2.3 and Remark 2.5, together with Axiom K, imply the axioms in Figure 1. Instead of doing that, in this section, we just check that the axioms are provable from inductive definitions of equality and dependent product types. One can make these inductive definitions in Agda as follows:

```
data _≡_ {l}{A : Set l} : {B : Set l} → A → B → Set l where
  rfl : (x : A) → x ≡ x
data ∑ {l m}{A : Set l}(B : A → Set m) : Set (l ⊔ m) where
  _,_ : (x : A) → B x → ∑ A B
-- the derived homogeneous equality
_≡_ : ∀{l}{A : Set l} → A → A → Set l
x ≡ y = x ≡ y
```

Then Agda's implementation of dependent pattern matching enables straightforward definitions of the functions from Figure 1, as follows:

```
ctr : ∀{l}{A B : Set l}{x : A}{y : B}(e : x ≡ y) → rfl x ≡ e
ctr (rfl x) = rfl (rfl x)

eqt : ∀{l}{A B : Set l}{x : A}{y : B} → x ≡ y → A ≡ B
eqt {_,_} {A} (rfl _) = rfl A

tpt : ∀{l m n}{A : Set l}{B : A → Set m}(C : (x : A) → B x → Set n)
      {x x' : A}{y : B x}{y' : B x'} → x ≡ x' → y ≡ y' → C x y → C x' y'
tpt _ (rfl _) (rfl _) y = y

module _ {l m}{A : Set l}{B : A → Set m} where
  fst : ∑ A B → A
  fst (x , _) = x

  snd : (z : ∑ A B) → B (fst z)
  snd (_, y) = y

  fpr : (x : A)(y : B x) → fst (x , y) ≡ x
  fpr x _ = rfl x

  spr : (x : A)(y : B x) → snd (x , y) ≡ y
  spr _ y = rfl y

  eta : (z : ∑ A B) → (fst z , snd z) ≡ z
  eta (x , y) = rfl (x , y)
```

Since we know from the previous section that these functions entail Axiom K, the above definitions have to use Agda's default `--with-K` option to switch the existing implementation of dependent pattern matching [6] back to the original version due to Reference [8], which is known

to imply Axiom K [10]. More precisely, it is only the matches on the two occurrences of the pattern `refl_` in the definition of `tpt` that involve an implicit use of Axiom K (to discharge the unification constraints $A \doteq A$ and $B\ x \doteq B\ x$); all the other functions can be defined without Axiom K.

4 CONCLUSION

This article has investigated heterogeneous equality and produced a simple collection of axioms for its typal form, in the spirit of Reference [9]. The point of view is foundational. From a practical perspective, the use of heterogeneous equality has much to recommend it for formalizing mathematics in dependent type theory when assuming uniqueness of identity proofs¹; but that is another story.

APPENDIX: TYPAL HOMOGENEOUS EQUALITY WITHOUT K

In this Appendix, for the sake of completeness, we consider axioms in dependent-type theory without Axiom K for *homogeneous* equality types

`postulate`

```
_≡_ : ∀{l}{A : Set l} → A → A → Set l
```

following Reference [9]. (Since without Axiom K heterogeneous equality is not very useful, we do not bother to consider axiomatizing `≡≡` in that setting.) One of the axioms makes use of dependent product types. Although one could axiomatize those types as we did in the main part of the article, it is simpler to use an inductive definition and corresponding pair patterns:

```
data ∑ {l m}{A : Set l}(B : A → Set m) : Set (l ⊔ m) where
  _,_ : (x : A) → B x → ∑ A B
- concrete syntax for ∑-types
syntax ∑ A (λ x → B) = ∑ x : A , B
- dependent product projections
module _ {l m}{A : Set l}{B : A → Set m} where
  fst : ∑ A B → A
  fst (x , _ ) = x
  snd : (z : ∑ A B) → B (fst z)
  snd (_ , y) = y
```

The axioms for homogeneous equality are

`postulate`

```
refl : ∀{l}{A : Set l}(x : A) → x ≡ x
```

```
cntr : ∀{l}{A : Set l}{x y : A}(e : x ≡ y) → (x , refl x) ≡ (y , e)
```

```
sbst : ∀{l m}{A : Set l}(B : A → Set m){x x' : A} → x ≡ x' → B x → B x'
```

Coquand also considers a regularity axiom for `sbst` (ax₃ in *loc.cit.*), but one can do without that by using Peter Lumsdaine's trick to correct `sbst` to a version `subst` for which there is a proof

¹Such is the approach of Lean [16] since version 3, for example.

`substIsRegular` : $\forall b \rightarrow \text{subst} (\text{refl } x) b \equiv b$, as follows. The proof begins as for Lemma 2.1 by considering functions that are injective modulo $\equiv$:

```
Inj :  $\forall \{l\} (A B : \text{Set } l) \rightarrow \text{Set } l$ 
Inj A B =  $\sum f : (A \rightarrow B) , \forall \{x y\} \rightarrow f x \equiv f y \rightarrow x \equiv y$ 

id :  $\forall \{l\} \{A : \text{Set } l\} \rightarrow A \rightarrow A$ 
id x = x

idInj :  $\forall \{l\} (A : \text{Set } l) \rightarrow \text{Inj } A A$ 
idInj _ = (id , id)
```

But then to construct `subst` and `substIsRegular`, one has to work a bit harder than in the proof of the lemma because of the lack of uniqueness of identity proofs:

```
module _ {l m} {A : Set l} (B : A  $\rightarrow$  Set m) {x : A} where
  Inj2 : {y z : A}  $\rightarrow x \equiv y \rightarrow x \equiv z \rightarrow \text{Inj } (B y) (B z)$ 
  Inj2 {y} p q =
    subst ( $\lambda z' \rightarrow \text{Inj } (B y) (B z')$ ) q
    (subst ( $\lambda y' \rightarrow \text{Inj } (B y') (B x)$ ) p (idInj (B x)))

  sbst2 : {y z : A}  $\rightarrow x \equiv y \rightarrow x \equiv z \rightarrow B y \rightarrow B z$ 
  sbst2 p q = fst (Inj2 p q)

C : {y : A} (p : x  $\equiv$  y) (b : B x)  $\rightarrow \sum c : B y , (\text{sbst}_2 p p c \equiv \text{sbst}_2 (\text{refl } x) p b)$ 
C p b = sbst C' (cntr p) (b , refl _)
  where
    C' :  $\sum y : A , (x \equiv y) \rightarrow \text{Set } m$ 
    C' (y , p) =  $\sum c : B y , (\text{sbst}_2 p p c \equiv \text{sbst}_2 (\text{refl } x) p b)$ 

subst : {y : A}  $\rightarrow x \equiv y \rightarrow B x \rightarrow B y$ 
subst p b = fst (C p b)

substIsRegular : (b : B x)  $\rightarrow \text{subst} (\text{refl } x) b \equiv b$ 
substIsRegular b = snd (Inj2 (refl x) (refl x)) (snd (C (refl x) b))
```

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